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Micro-Concurrent Pascal: 1802 System Evaluation

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Introduction

Micro-Concurrent Pascal, a Pascal dialect, is a high-level language that supports real-time, inter-tasking programming. The mCP language offers process and monitor constructs that allow any number of processes to run independently but at the same time share data and communicate with each other. An mCP program requires a certain amount of organization in defining the structure of these software constructs. This underlying structure, different from a standard Pascal program, promotes a systematic approach to creating error-free real time systems.

MCP programs are compiled on a host computer which performs extensive compile-time checking. The compiler produces pseudo code which may then be downloaded to the target system. On the target system a 4.6 kilobyte Interpreter/Kernel executes the program. The P-code (pseudo code) that is produced is position-independent, reentrant, and ROMable. The mCP program can also access assembly language routines.

The mCP program which must be interpreted will not execute as fast as an equivalent assembly language program. The 1802 Interpreter/Kernel will execute approximately 1950 P-codes per second (using a 2.4576 Mhz crystal). This P-code execution speed plays a major factor in the interrupt handling efficiency of mCP. Since most real time applications involve interrupts, mCP must be able to handle interrupts efficiently. This report will examine the specifics of the interrupt handling scheme of mCP. The Enertec real time clock program, ERTC18, will be used in order to outline the steps taken by the kernel in servicing an interrupt.

It is assumed that the reader is familiar with the mCP inter-processing organization. An outline of the ERTC18 program will be given although a more comprehensive description is given in the Enertec manual.

Steps in Handling Interrupts

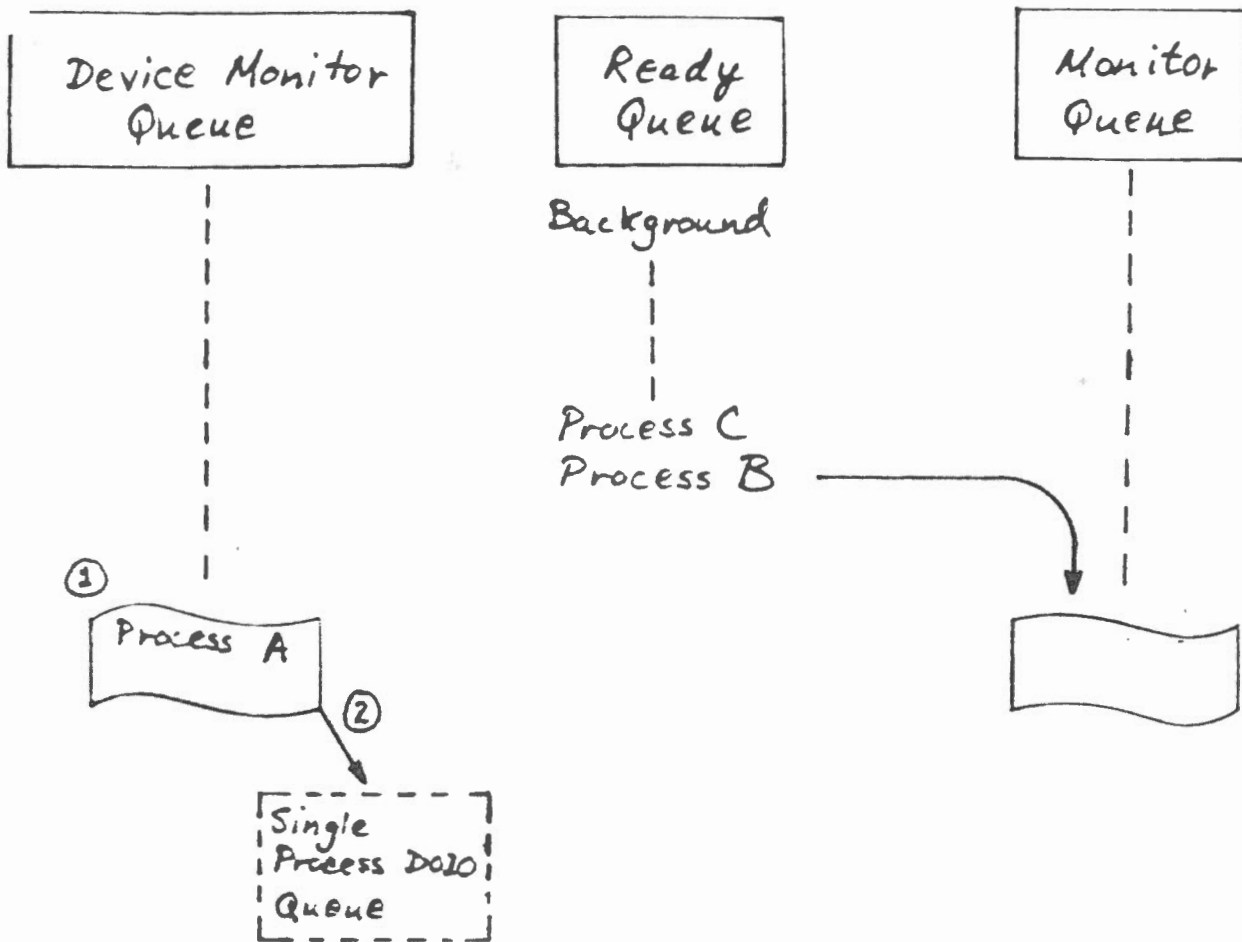
The interrupt handling system in mCP is based on the mCP DOIO P-code instruction. This instruction may only occur in a device monitor. On execution, it causes the currently running process to be preempted leaving that process to wait for an interrupt to occur. The kernel will then fetch the next process on the ready queue and allow it to run. When an interrupt occurs, the kernel identifies the source of interrupt, preempts the current running process, and resumes operation of the process that had been preempted and was waiting for the interrupt (see fig. 1).

In the present system up to 32 levels of interrupts can be handled, each having a unique priority, and these priorities may be altered dynamically.

The ERTC18 program operates a real time clock. The operator may set the current time from a terminal and view the current time on the terminal display. The program updates the time on receiving a 1 hertz interrupt from a clock source. In total, three interrupts must be recognized by the operating system: clock interrupt, UART receive and UART transmit interrupts, with the clock interrupt having the highest priority.

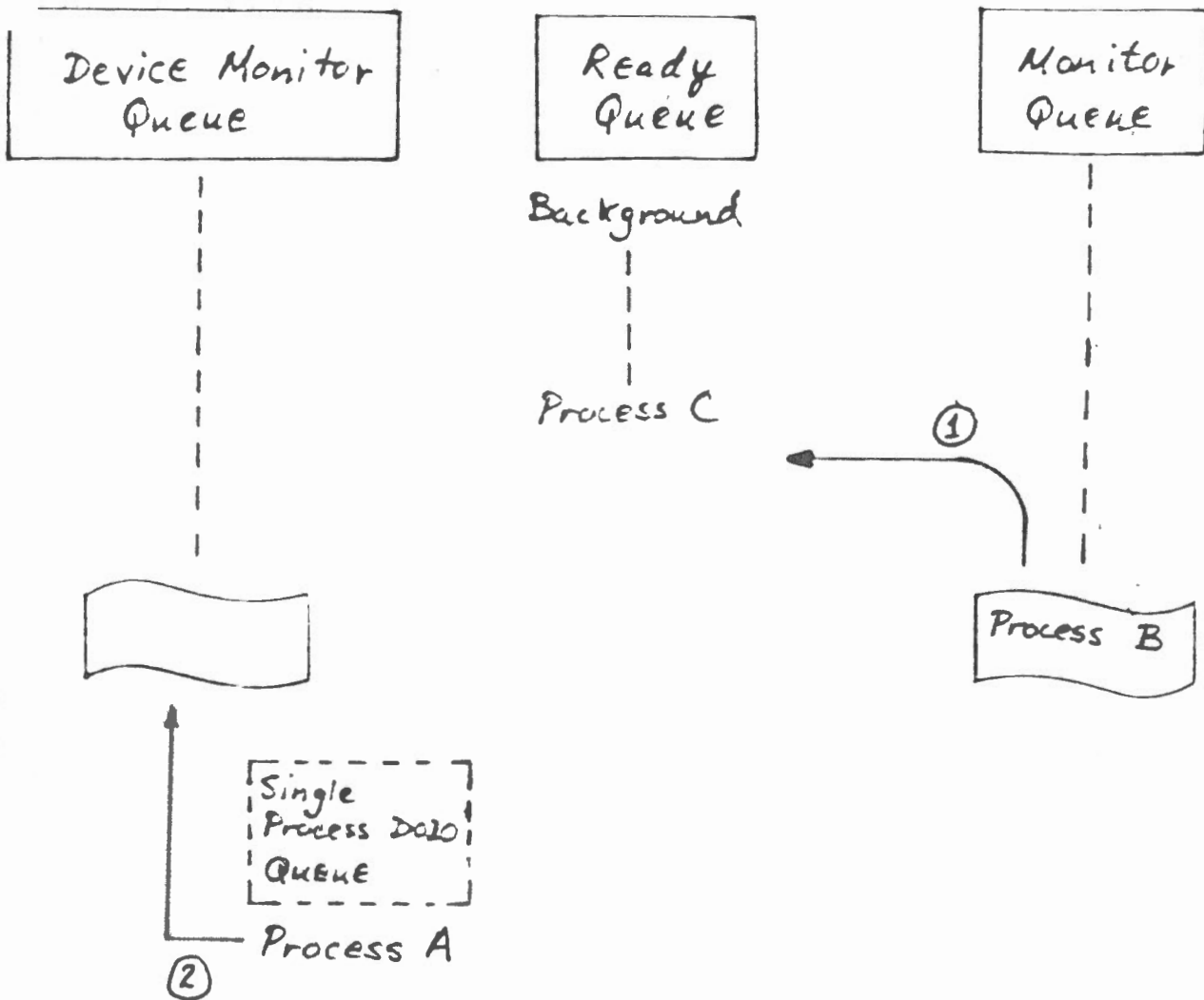
The program consists of two processes: the clock process and the operator process. These processes are continuous loops. The clock process handles the clock interrupts in the clock device monitor, and encapsulates the time data structure in the clock monitor. The operator process allows interaction between the user and the system. It accesses the Uart read and write device monitors and allows the user to access the time through the clock monitor. The clock process is continuously running, responding to clock interrupts and updating the time data in the clock monitor. The clock monitor can be accessed by the operator process and thus serves as the bridge between the two processes. The operator process continually waits for operator input to read or change the current time in the clock monitor.

FIGURE 1.(a) Simple System Response to a DOIIO Instruction



1. Process A active in Device Monitor issues a DOIIO statement.
2. Process A preempted and placed on the device monitor's Single Process DOIIO Queue.
3. Process at head of Ready Queue is run (i.e. Process B enters Monitor).

FIGURE 1. (b)



1. Interrupt occurs, currently running process (Process B) is preempted and placed at head of Ready Queue.
2. Process waiting for interrupt, Process A, restarted (in device monitor).

The program flow is shown in figure 2. When the program is run, the time is reset to 00:00:00. In the clock device monitor a DOIO statement preempts the clock process and the clock process then waits for an interrupt. The other process, the operator process, is then scheduled, initializing the UART and sending a prompt (">") to the terminal. The operator process is then preempted on execution of a DOIO statement in the UART read device monitor, and the process then waits for the operator to enter a command on the terminal.

At this point the only process which is running is the background process. Both clock and operator processes have been preempted and are placed on the clock single process DOIO queue and the UART single process DOIO queue, respectively. The system in this state will from now on be referred to as the "minimum state". It is in the minimum state because all monitor queues are empty and the only process on the ready queue is the background process (see figure 3).

FIGURE 2 ERTC18 Program Flow

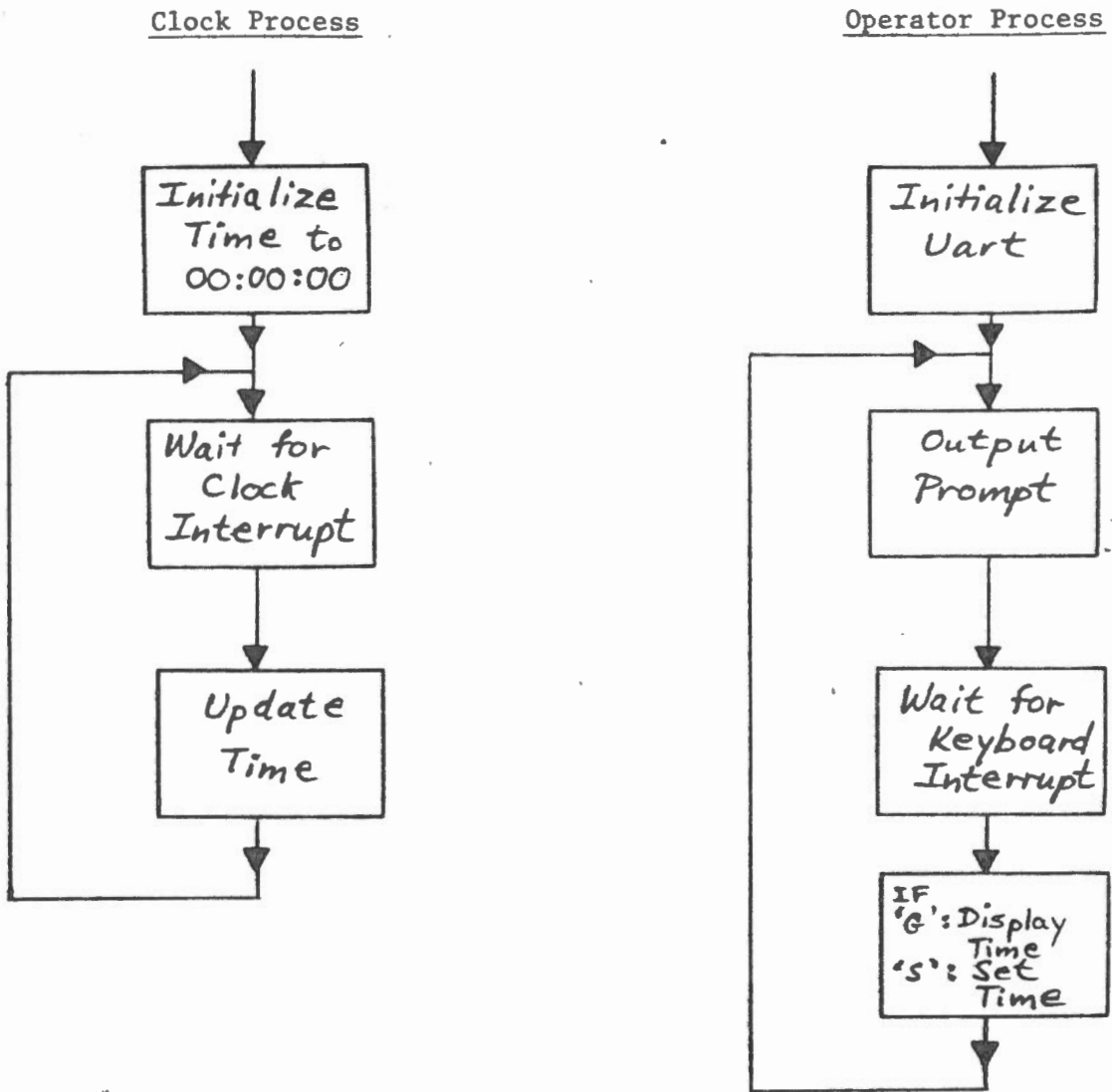


Figure 3. Minimum State of ERTC18 Operating System

READY
QUEUE

Background
Process

CLOCK
MONITOR
QUEUE

CLOCK
DEVICE
MONITOR
QUEUE

Single Process
DIO QUEUE

Clock Process

UART READ
DEVICE
MONITOR
QUEUE

Single Process
DIO QUEUE

Operator Process

UART WRITE
DEVICE
MONITOR
QUEUE

Single Process
DIO QUEUE

Clock Interrupt Servicing

During the execution of a process in mCP the kernel will disable interrupts as a P-code instruction is fetched, and enable interrupts after a P-code instruction is completed. Hence there is a "window" through which interrupts are allowed to interrupt the running process. An interrupt could occur just after a P-code was fetched, and the kernel would then prevent the CPU from being interrupted until the P-code was interpreted and executed. This reduces the efficiency of interrupt handling. Although this problem cannot be directly avoided, it points to the major contributing factor in the efficiency of interrupt servicing: the speed of executing P-codes.

In the minimum state of the ERTC18 program, the running process is the background process. This process consists of a mCP HALT instruction which puts the CPU in an idle state and waits for interrupts. Thus the problem of an interrupt occurring while the kernel is executing a P-code is avoided. As soon as any interrupt occurs in the minimum state, control is transferred immediately to the kernel's interrupt handling routines. The time taken by the kernel to handle interrupts from this state will be the minimum.

The code contained in the Clock device monitor must be analyzed to determine the interrupt handling efficiency. This code is listed below,

```
BEGIN
        DOIO;
        OUT( #00,CLOCK_WORD );
END
```

The DOIO statement suspends the clock process to wait for an interrupt to occur. The operator process is run, and eventually suspended by a DOIO statement. The system is in the minimum state with both processes waiting for interrupts. When the clock interrupt occurs the Kernel's interrupt servicing follows these steps:

- 1: The source of Interrupt is determined.
- 2: The presently running process (the background process) is preempted.
- 3: The process waiting for the interrupt (i.e. clock process) is restarted.

The source of interrupt is determined by examining the event flag associated with the interrupting device (see Appendix: Clock Hardware). When the clock process is restarted, the next P-code fetched and executed will be the OUT instruction. This instruction will cause a pulse to be transmitted which will reset the clock hardware. The clock device monitor will then be exited on execution of the END statement. The clock process then updates the software time variable by one (i.e. one second) and eventually the process cycles back to the minimum state upon execution of the DOIO instruction in the clock device monitor. Every second, on clock interrupts, this sequence is repeated.

The steps in handling the clock interrupt from the minimum state are outlined in figure 4. Labels identify the routines used in the kernel. From measurements made on the logic analyzer, it was found that it takes 238 microseconds for the kernel to recognize the highest priority Legal interrupt (i.e. the clock interrupt, Step 1).

At this point additional time is required to preempt the current running process and switch control to the process waiting for the interrupt. The time it takes between switching processes up to the point of fetching the OUT P-code in the device monitor following the DOIO instruction is 694 microseconds (Steps 2 and 3).

Execution of the OUT instruction to reset the clock takes 2621 microseconds. Up to this point the total time taken from the time the clock interrupt occurred to the time when the clock hardware is reset is 3552 microseconds.

Finally, the time taken for the kernel to exit from the device monitor and return to the clock process (where the time variable is to be updated) is 4493 microseconds.

FIGURE 4.(a) Kernel Steps in Handling Interrupts

INTVEC

TRYME

LEGAL

NEWSTA

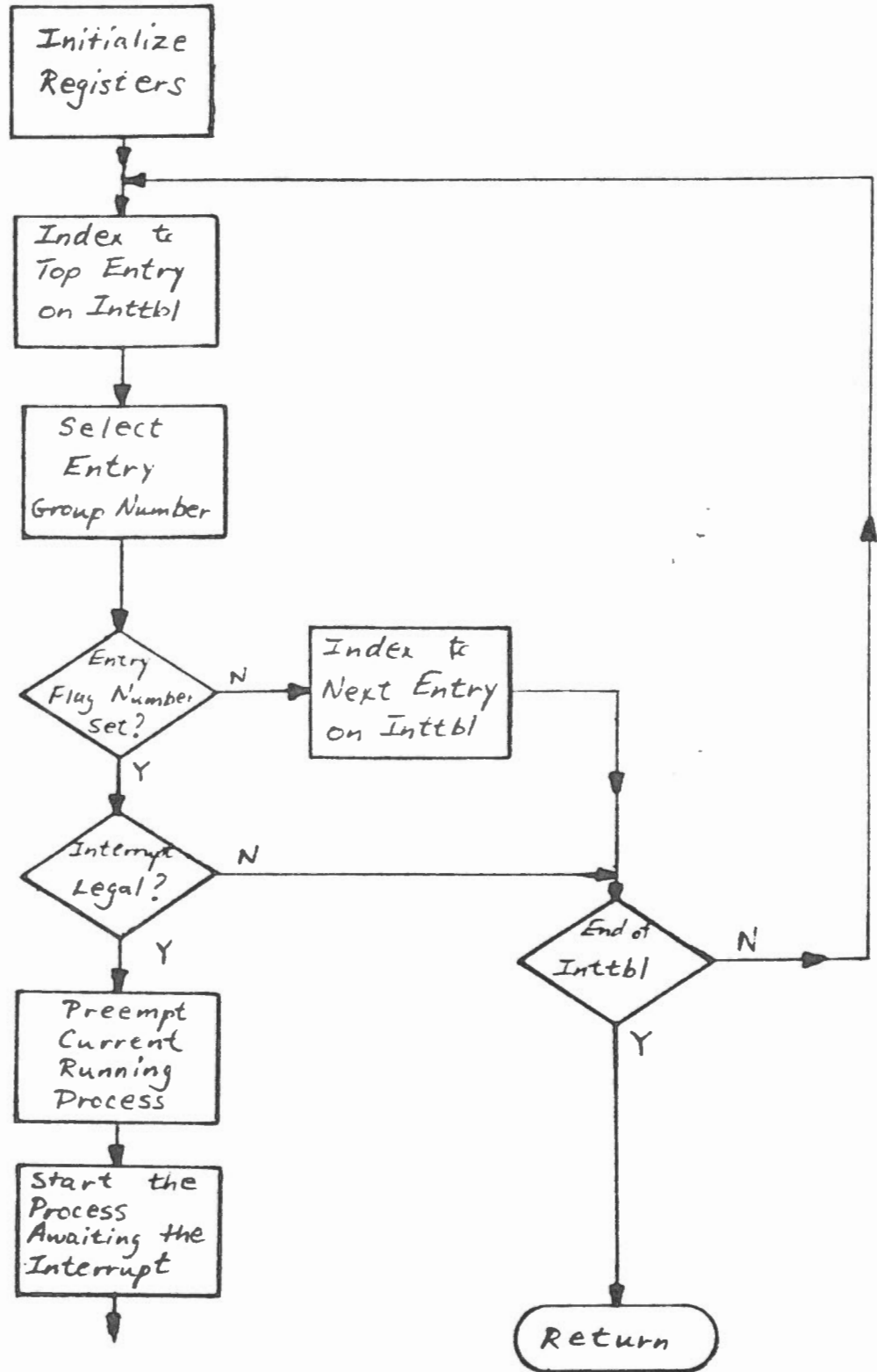
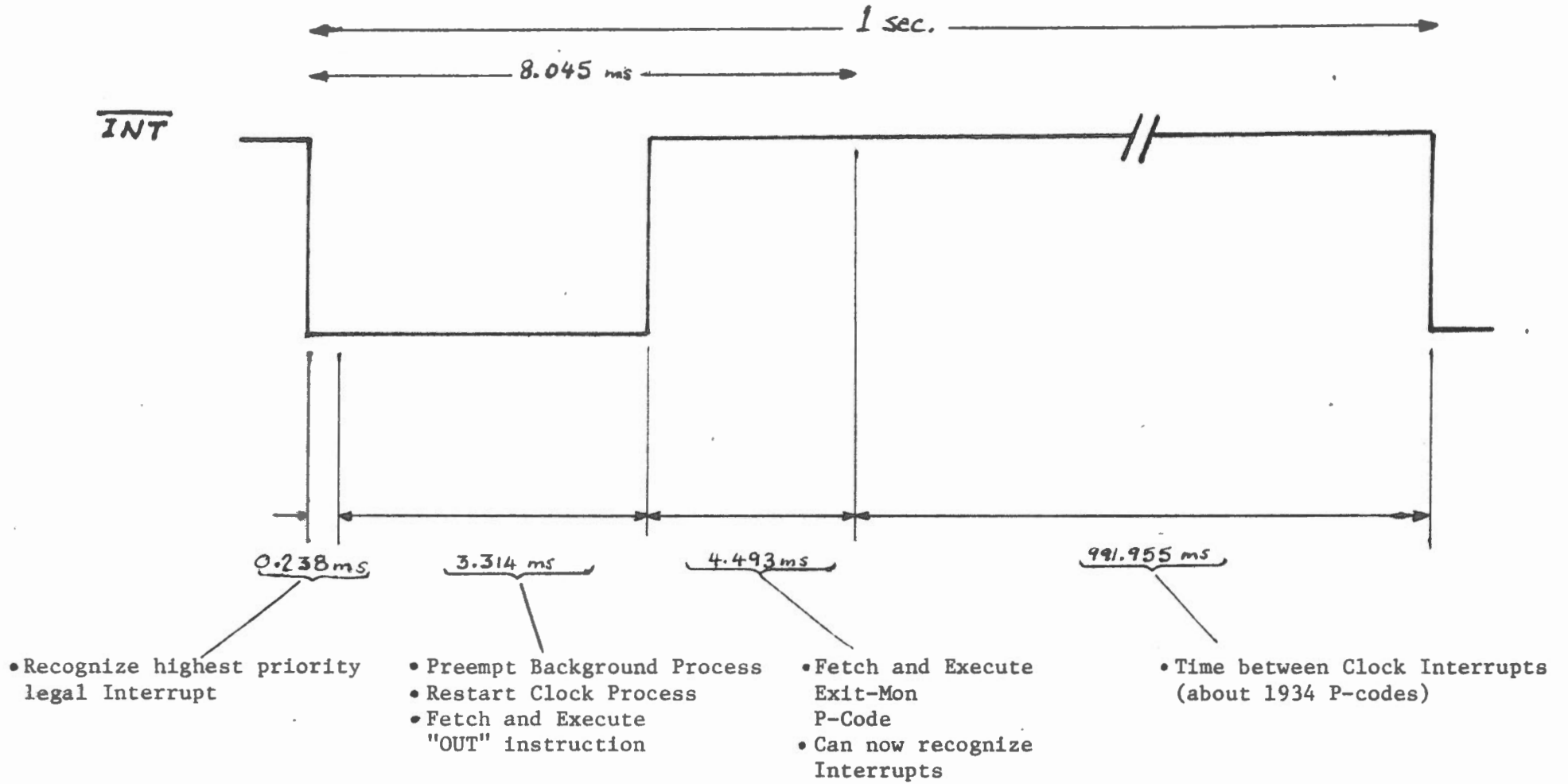


Figure 4.(b) Clock Interrupt Servicing



MCP Interrupt Handling Problems: Summary

The organization of the ERTC18 program is satisfactory for the intended purpose - providing a simple operating system allowing a real time clock to be set or viewed via a terminal. The system can handle the 1 Hz clock interrupt effectively. Since the interrupt period is sufficiently long there is enough time after the clock interrupt has been handled for the system to execute up to 1934 P-codes. However, if we wanted to run the clock at higher speeds problems would arise.

A clock frequency of 200 Hz for example could not be handled since it takes 8.045 milliseconds to recognize the interrupt, reset the clock, and exit from the clock device monitor. A clock frequency of 120 Hz could be handled, but there would not remain any time between interrupts to run other processes. At 5 Hz, though, there would be 192 milliseconds between interrupts, and 374 P-codes could be executed. This clock frequency would be the most appropriate for the present 1802 system.

Other problems will arise in multiple interrupting applications. MCP does not enable interrupts after I/O instructions. In the ERTC18 program although the clock interrupt is recognized and reset in 3.55 milliseconds, the kernel will only enable interrupts when the P-code following the OUT P-code has been completed. This means that the CPU is tied up for another 391 microseconds where it cannot service interrupts from other sources. Furthermore, the time taken to exit from the device monitor that resets the interrupting device and switch processes back to the main running process will reduce the available time to perform background processing (i.e. data manipulations, math operations, etc.).

For minimum-processing applications like the ERTC18 program where the frequency and number of interrupts are low, then the present MCP interrupt handling scheme is acceptable. A clock interrupt frequency of 5 kHz would allow enough time between interrupt servicing for 374 P-codes, for example. For applications requiring higher clock interrupt frequencies, with several sources of interrupts, and more processing, the MCP interrupt handling scheme will not work effectively and may have to be altered.

Hardware/Software Solutions

The simplest way to increase the clock interrupt handling speed is to remove the mCP OUT instruction that resets the clock hardware. Knowing that it takes 238 microseconds to recognize a legal interrupt, we could reduce the clock interrupt pulse to 500 microseconds. This would allow sufficient time for the kernel to identify the interrupt while ensuring that the interrupt line will be reset before the clock device monitor is exited. By removing the OUT instruction a saving of 2621 microseconds can be achieved. The total time taken to handle this interrupt, from the moment the interrupt occurs to the time interrupts are enabled after exiting from the device monitor will be 5.424 ms compared to the previous 8.045 ms.

An alternate safer approach would require modifying the kernel. When the appropriate interrupt is identified, a "Set Q", "Reset Q" sequence could be performed within the kernel, thus sending a pulse via the CPU Q output. This output would reset the clock hardware before the next P-code (i.e. EXIT) was fetched. This would replace the OUT instruction, and thus the time savings would be similar to the previous approach.

These methods are simple, yet would not provide any real great improvement in general interrupt handling. At most an additional 5 P-codes could be executed between the clock interrupts. Any need for faster handling would require modifications to the basic mCP program design, necessitating kernel modifications and a departure from the DOIO "wait-for-interrupt" approach. Some ideas include interrupt vectoring or a dedicated clock interrupt service routine.

Interrupt vectoring would require some external hardware such as the RCA 1877 Interrupt Controller. Modifications to the mCP kernel and the DOIO section of the interpreter would have to be made.

The other possibility is in applications where a clock interrupt plays a major role in the scheduling of processes and control. Use of the RCA 1804 CPU with its internal counter-timer could provide accurate presettable internal clock interrupts. Modifications of the kernel would allow conditional Branch on Internal Interrupts to transfer control to certain processes or simply update a software counter.

Conclusion

The aim of this report was to evaluate the 1802 micro-Concurrent Pascal system. After running the Enertec ERTC18 real-time program it was found that the clock interrupt servicing was not very efficient. This report examined the MCP interrupt handling scheme particular to the 1802 system, using the ERTC18 program as an example.

An analysis of the steps taken by the MCP kernel in servicing interrupts showed that for the ERTC18 program, handling of the clock interrupt was from a minimum state. From this state, CPU servicing of the interrupt was the quickest possible. The time taken to recognize this interrupt, reset the clock interrupt device, and exit from the clock device monitor was 8.045 milliseconds.

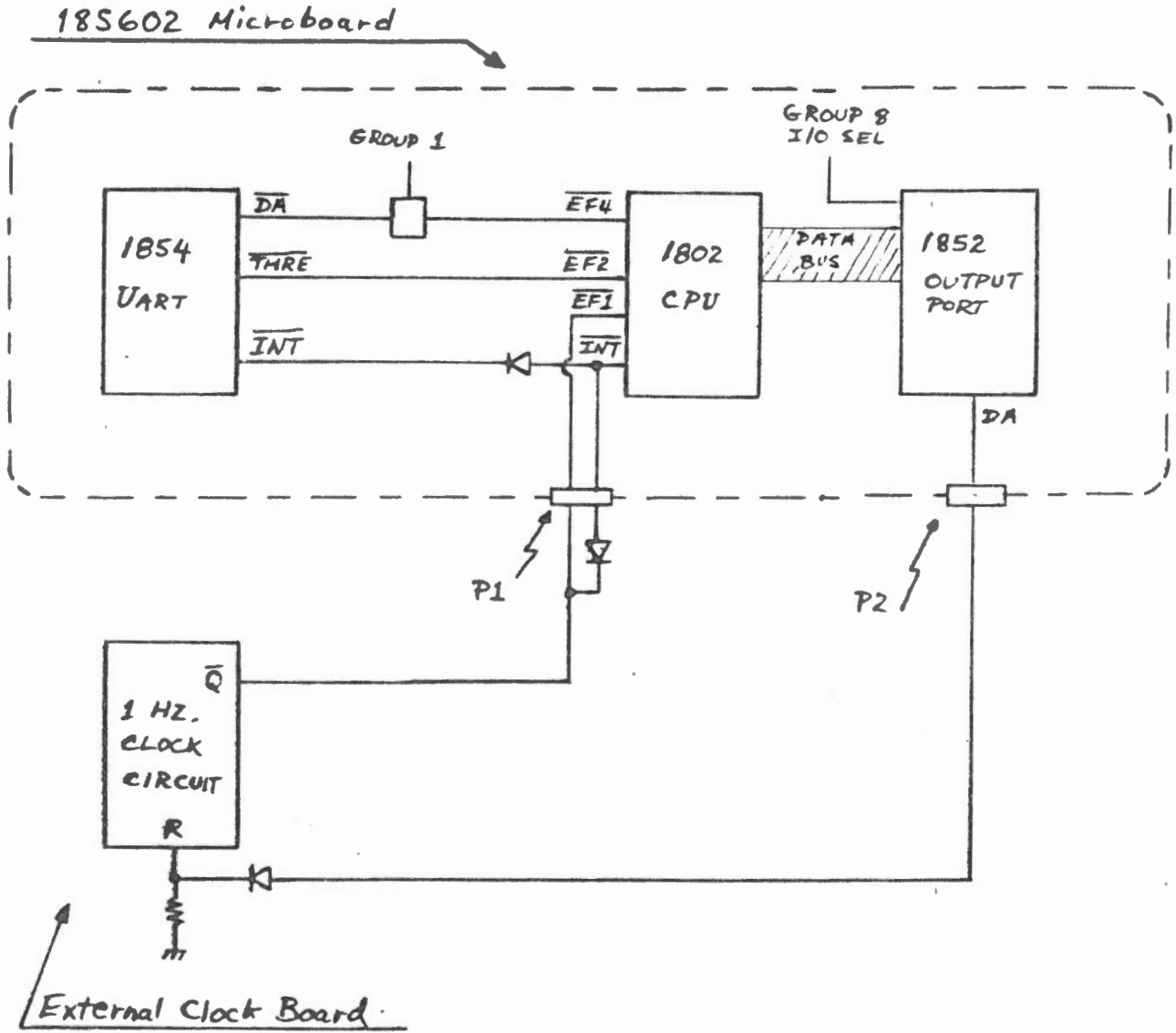
For applications such as the ERTC18 real-time clock, where the frequency and number of interrupting devices is low, the present MCP system is suitable. The time available between the 1 hertz clock interrupt was enough to allow for updating of the data and examining the UART device for operator input. A suitable typical application could be a data acquisition system that requires a terminal control and performs measurements on a 5 Hz clock interrupt. In this case there would be enough time for about 374 P-codes to be executed between clock interrupt servicing. If the data processing overhead is not too severe then the present system will perform well for this type of application.

Applications involving several, frequent interrupting devices or heavy mathematical computation and data processing between interrupts, then the present system will not work well. In such a system requiring quick response to an interrupt, certain hardware or software modifications will have to be made. A modification of the kernel could allow a faster acknowledgement of a clock interrupt and avoid the necessity of having a clock device monitor and its software overhead to control the clock. MCP's provision for accessing external assembly language routines can be used advantageously, or hardware controlling of interrupts could benefit certain higher speed applications.

In general, though, these modifications would require a modification of the MCP interrupt handling scheme. This would entail a departure from the structured, portable, high-level software to a more dedicated, less portable system. With an interpretive pseudo-coded language such as MCP, code execution speed will not be as fast as a similar assembly language program. The fundamental problem of interrupt handling for real-time applications lies not in the MCP interrupt servicing scheme, but in the speed of P-code execution. Attempts to modify the kernel can improve a system, but the overall performance of the system is limited by the speed of P-code execution. Modifications of the system should be incorporated within the basic MCP programming constructs (i.e. processes, monitors, and device monitors). In this way the fundamental reasons for MCP programming can be realized. A real-time program can be created using the clarity and flexibility of a structured high-level language with a programming environment that reduces the possibility of real-time errors.

Appendix 1

System Device Interrupt Connections



<u>Interrupt Signal</u>	<u>Group Number</u>	<u>Associated Event Flag</u>
1 Hz Clock	8	$\overline{EF1}$
\overline{DA}	1	$\overline{EF4}$
\overline{THRE}	None	$\overline{EF2}$

Appendix 2

```

($LIST=LONG, DERUG ON                                     $)
(#####)
#
#   Micro Concurrent Pascal (MCP) Program                 #
#   ENERTEC Real-Time for the 1802                       #
#   4/13/81                                              #
#
#   (c) Copyright   1981                                 #
#
#   ENERTEC, Inc.                                       #
#   19 Jenkins Avenue                                  #
#   Lansdale, Pa. 19446                                #
#   (215) 362-0966                                     #
#
#####)

```

```

CONST CLOCK_GROUP = 8;
      CLOCK_FLAG = 1;
      CLOCK_SELECTOR = 6;
      UART_GROUP = 1;
      UART_RECV_FLAG = 4;
      UART_RECV_SELECTOR = 4;
      UART_XMIT_FLAG = 2;
      UART_XMIT_SELECTOR = 3;
      INTTBL_TERMINATOR = -1;

```

```

STRUC_CON INTTBL: ARRAY[1..10] OF -1..8 =
  [CLOCK_GROUP, CLOCK_FLAG, CLOCK_SELECTOR,
   UART_GROUP, UART_RECV_FLAG, UART_RECV_SELECTOR,
   UART_GROUP, UART_XMIT_FLAG, UART_XMIT_SELECTOR,
   INTTBL_TERMINATOR];

```

```

CONST DATA_WORD = ADR(#0102);
      CTRL_WORD = ADR(#0103);
      CLOCK_WORD = ADR(#0802);
      LINELENGTH = 74;
      NUL='(:0:)' ; BS='(:8:)' ; LF='(:10:)' ; CR='(:13:)' ;
      NUL2='(:0:)(:0:)' ;

```

```

TYPE TIME = RECORD
      HOUR : 0..24;
      MIN, SEC : 0..60
    END;
      LINE_DISP =(PROMPT, NEWLINE, STAY);
      LINE = ARRAY[1..LINELENGTH] OF CHAR;
      INT = 0..127;

```

```

(#####)
#   UART_WRITE #
(#####)

```

```

TYPE UART_WRITE = DEVICE_MON (SELECTOR: INT);

```

```

PROCEDURE ENTRY WRITE(MESSAGE: LINE; DISP: LINE_DISP);
  VAR I: INT;
      THROWAWAY: INTEGER;
  BEGIN

```

```

    I:=1;
    OUT(#BD, CTRL_WORD) ( XMIT REQ., INT. EN., 8 DATA, 2 STOP, NO PARITY
    DOIO;
    WHILE (MESSAGE[I] <> NUL) AND (I<LINELENGTH) DO
    BEGIN
      OUT(ORD(MESSAGE[I]) DATA_WORD); (SEND A CHARACTER)

```

```

        INC(I);
    END;
    IF (DISP=PROMPT) OR (DISP=NEWLINE) THEN
        BEGIN OUT(ORD(CR), DATA_WORD); DOIO;
              OUT(ORD(LF), DATA_WORD); DOIO;
        END;
    IF DISP=PROMPT THEN
        BEGIN OUT(ORD(')'), DATA_WORD); DOIO; END;
        OUT(#3D, CTRL_WORD); {TRANSMIT INHIBIT OTHERWISE SAME AS ABOVE}
        THROWAWAY := INN(CTRL_WORD);
    END;

```

```

PROCEDURE ENTRY ECHO(CHRS:CHAR);
    BEGIN
        OUT(#BD, CTRL_WORD);
        OUT(ORD(CHRS), DATA_WORD);
        OUT(#3D, CTRL_WORD);
    END;

```

```

    BEGIN
        OUT(#3D, CTRL_WORD);
    END;

```

```

(#####
# UART_READ #
#####)

```

```

TYPE UART_READ = DEVICE_MON (TERM_OUT: UART_WRITE;
                              SELECTOR: INT);

```

```

    VAR THROWAWAY: INT;

```

```

PROCEDURE ENTRY READ( VAR MESSAGE: LINE; VAR LENGTH: INT);
    VAR LASTCHAR: CHAR;
    BEGIN
        LENGTH:=0;
        REPEAT
            IF LENGTH<LINELENGTH THEN INC(LENGTH);
            DOIO;
            LASTCHAR := CHR(INN(DATA_WORD)); {GET A CHARACTER}
            IF LASTCHAR=BS THEN
                BEGIN
                    IF LENGTH >= 2 THEN DEC(LENGTH);
                    DEC(LENGTH);
                END
            ELSE MESSAGE[LENGTH] := LASTCHAR;
            IF LASTCHAR (<) CR THEN
                BEGIN
                    TERM_OUT.ECHO(LASTCHAR);
                END;
            UNTIL (LASTCHAR = CR);
            TERM_OUT.WRITE(NUL2, NEWLINE);
        END;

```

```

    BEGIN
        THROWAWAY:=INN(DATA_WORD);
        THROWAWAY:=INN(DATA_WORD);
    END;

```

```

(#####
# CLOCK_MONITOR #
#####)

```

```

TYPE CLOCK_MONITOR = MONITOR (TERM_OUT: UART_WRITE);

```

```

PROCEDURE ENTRY TICK;
  BEGIN
    WITH CLOCKTIME DO
      BEGIN
        INC(SEC);
        IF SEC=60 THEN
          BEGIN
            SEC:=0; INC(MIN);
            IF MIN=60 THEN
              BEGIN
                MIN:=0; INC(HOUR);
                IF HOUR=24 THEN
                  BEGIN
                    HOUR:=0;
                    TERM_OUT.WRITE('IT IS A NEW DAY(::0:)', NEWLINE);
                  END;
                END
              END
            END
          END
        END;
      END;
    END;
  END;

```

```

PROCEDURE ENTRY SETTIME(T: TIME);
  BEGIN
    CLOCKTIME := T;
  END;

```

```

PROCEDURE ENTRY GETTIME(VAR T: TIME);
  BEGIN
    T := CLOCKTIME;
  END;

```

```

BEGIN
  WITH CLOCKTIME DO BEGIN HOUR:=0; MIN:=0; SEC:=0; END;
END;

```

```

(*****
# CLOCKPULSE #
*****)

```

```

TYPE CLOCKPULSE = DEVICE_MON (SELECTOR: INT);

```

```

PROCEDURE ENTRY TICK;
  BEGIN
    DOIO;
    OUT(#00, CLOCK_WORD);
  END;

```

```

BEGIN
  ( SET UP TO START THE CLOCK )
END;

```

```

(*****
# CLOCKPROCESS #
*****)

```

```

TYPE CLOCKPROCESS = PROCESS (PULSE: CLOCKPULSE;
                              CLOCK: CLOCK_MONITOR);

```

```

BEGIN
  CYCLE CLOCK.TICK; PULSE.TICK; END;
END;

```

```

(*****
# OPERATORPROCESS #
*****)

```

```
TYPE OPERATORPROCESS = PROCESS ( CLOCK: CLOCK_MONITOR;  
                                TERM_IN: UART_READ;  
                                TERM_OUT: UART_WRITE);
```

```
VAR BUFFER: LINE;  
    LENGTH: INT;  
    T: TIME;
```

```
FUNCTION NUMBER(CCHARACTER: CHAR): INT;  
    BEGIN NUMBER:=ORD(CCHARACTER) - ORD('0'); END;
```

```
FUNCTION ASCII(NUMBER: INT): CHAR;  
    BEGIN ASCII:=CHR(NUMBER + ORD('0')); END;
```

```
BEGIN
```

```
    CYCLE
```

```
        TERM_OUT.WRITE(NUL2, PROMPT);
```

```
        TERM_IN.READ(BUFFER, LENGTH);
```

```
        CASE BUFFER[1] OF
```

```
'S' (SET TIME) :
```

```
    BEGIN
```

```
        WITH T DO
```

```
            BEGIN
```

```
                HOUR:=23;
```

```
                MIN :=50;
```

```
                SEC :=10;
```

```
            END;
```

```
        CLOCK.SETTIME(T);
```

```
    END;
```

```
'G' (GET TIME) :
```

```
    BEGIN
```

```
        CLOCK.GETTIME(T);
```

```
        WITH T DO
```

```
            BEGIN
```

```
                BUFFER[1]:=ASCII(HOUR DIV 10); BUFFER[2]:=ASCII(HOUR MOD 10);
```

```
                BUFFER[3]:=':';
```

```
                BUFFER[4]:=ASCII(MIN DIV 10); BUFFER[5]:=ASCII(MIN MOD 10);
```

```
                BUFFER[6]:=':';
```

```
                BUFFER[7]:=ASCII(SEC DIV 10); BUFFER[8]:=ASCII(SEC MOD 10);
```

```
                BUFFER[9]:=NUL;
```

```
            END;
```

```
        TERM_OUT.WRITE(BUFFER, NEWLINE);
```

```
    END;
```

```
(<): BEGIN END;
```

```
    END;
```

```
END;
```

```
END;
```

```
(*****
```

```
# INITIAL PROCESS #
```

```
*****)
```

```
VAR CLOCK: CLOCK_MONITOR;
```

```
    TICKTOCK: CLOCKPROCESS;
```

```
    PULSE: CLOCKPULSE;
```

```
    OPERATOR: OPERATORPROCESS;
```

```
    TERM_IN: UART_READ;
```

```
    TERM_OUT: UART_WRITE;
```

```
(MONITOR)
```

```
(PROCESS TO CYCLE CLOCK)
```

```
(DEVICE MONITOR FOR TIMER)
```

```
(PROCESS)
```

```
(DEVICE MON FOR TERMINAL)
```

```
(DEVICE MON FOR TERMINAL)
```

```
BEGIN
```

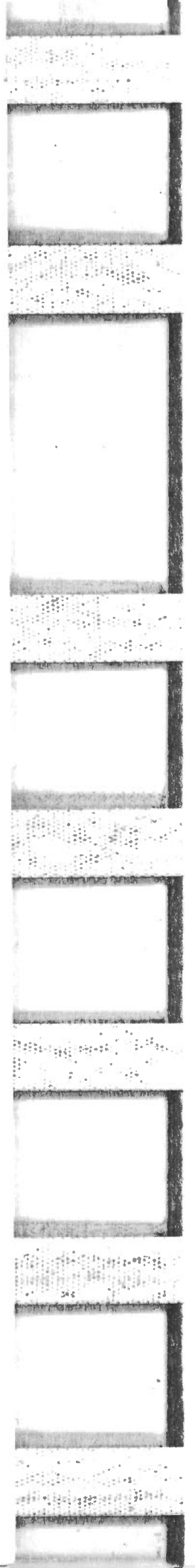
```
    INIT TERM_OUT(UART_XMIT_SELECTOR);
```

```
    INIT CLOCK(TERM_OUT);
```

```
    INIT PULSE(CLOCK_SELECTOR);
```

```
    INIT TERM_IN(TERM_OUT, UART_RECV_SELECTOR);
```


T TICKTOCK(PULSE, CLOCK);
T OPERATOR(CLOCK, TERM_IN, TERM_OUT);



PROGRAM LISTING

RLIN H

ELENGTH: 643
STLENGTH: 44

```
1: ( 0004)
2: ( 0004)      ($LIST=LONG, DEBUG ON                                $)
3: ( 0004)      {#####}
4: ( 0004)      #
5: ( 0004)      #   Micro Concurrent Pascal (MCP) Program          #
6: ( 0004)      #       ENERTEC Real-Time for the 1802            #
7: ( 0004)      #                               4/13/81              #
8: ( 0004)      #
9: ( 0004)      #           (c) Copyright   1981                    #
10: ( 0004)     #
11: ( 0004)     #           ENERTEC, Inc.                            #
12: ( 0004)     #           19 Jenkins Avenue                       #
13: ( 0004)     #           Lansdale, Pa. 19446                     #
14: ( 0004)     #           (215) 362-0966                          #
15: ( 0004)     #
16: ( 0004)     ######}
17: ( 0004)
18: ( 0004) CONST CLOCK_GROUP = 8;
19: ( 0004)     CLOCK_FLAG = 1;
20: ( 0004)     CLOCK_SELECTOR = 6;
21: ( 0004)     UART_GROUP = 1;
22: ( 0004)     UART_RECV_FLAG = 4;
23: ( 0004)     UART_RECV_SELECTOR = 4;
24: ( 0004)     UART_XMIT_FLAG = 2;
25: ( 0004)     UART_XMIT_SELECTOR = 3;
26: ( 0004)     INTTBL_TERMINATOR = -1;
27: ( 0004)
28: ( 0004) STRUC_CON INTTBL: ARRAY[1..10] OF -1..8 =
29: ( 0004)     [CLOCK_GROUP, CLOCK_FLAG, CLOCK_SELECTOR,
30: ( 0004)     UART_GROUP, UART_RECV_FLAG, UART_RECV_SELECTOR,
31: ( 0004)     UART_GROUP, UART_XMIT_FLAG, UART_XMIT_SELECTOR,
32: ( 0004)     INTTBL_TERMINATOR];
33: ( 0004)
34: ( 0004) CONST DATA_WORD = ADR(#0102);
35: ( 0004)     CTRL_WORD = ADR(#0103);
36: ( 0004)     CLOCK_WORD = ADR(#0802);
37: ( 0004)     LINELENGTH = 74;
38: ( 0004)     NUL='(:0:)' ; RS='(:8:)' ; LF='(:10:)' ; CR='(:13:)' ;
39: ( 0004)     NUL2='(:0:)(:0:)' ;
40: ( 0004)
41: ( 0004) TYPE TIME = RECORD
42: ( 0004)     HOUR ; 0..24;
43: ( 0004)     MIN, SEC ; 0..60
44: ( 0004)     END;
45: ( 0004)     LINE_DISP =(PROMPT, NEWLINE, STAY);
46: ( 0004)     LINE = ARRAY[1..LINELENGTH] OF CHAR;
47: ( 0004)     INT = 0..127;
48: ( 0004)
49: ( 0004) {#####}
50: ( 0004) # UART_WRITE #
51: ( 0004) #####}
```

```

END;
432E JUMP -210
END;
D6 END_PROCESS

```

```

(#####
# INITIAL PROCESS #
#####)

```

```

VAR CLOCK: CLOCK_MONITOR; (MONITOR)
    TICKTOCK: CLOCKPROCESS; (PROCESS TO CYCLE CLOCK)
    PULSE: CLOCKPULSE; (DEVICE MONITOR FOR TIMER)
    OPERATOR: OPERATORPROCESS; (PROCESS)
    TERM_IN: UART_READ; (DEVICE MON FOR TERMINAL)
    TERM_OUT: UART_WRITE; (DEVICE MON FOR TERMINAL)

```

```

) BEGIN
) INIT TERM_OUT(UART_XMIT_SELECTOR);
) BAF4 PS_AD_G_1 -12
) 03 CONSTANT 3
) C600000231FE INIT_MON 0 2 -463
) INIT CLOCK(TERM_OUT);
) BAFE PS_AD_G_1 -2
) B2F4 PS_VW_G_1 -12
) C6030002DAFE INIT_MON 3 2 -294
) INIT PULSE(CLOCK_SELECTOR);
) BAFA PS_AD_G_1 -6
) 06 CONSTANT 6
) C6000002EEFE INIT_MON 0 2 -274
) INIT TERM_IN(TERM_OUT, UART_RECV_SELECTOR);
) BAF6 PS_AD_G_1 -10
) B7 PS_VW_G_1 -12
) 04 CONSTANT 4
) C601000461FE INIT_MON 1 4 -415
) INIT TICKTOCK(PULSE, CLOCK);
) BAFC PS_AD_G_1 -4
) B2FA PS_VW_G_1 -6
) B2FE PS_VW_G_1 -2
) CB4100000004D9FE INIT_PROCESS 65 0 4 -295
) INIT OPERATOR(CLOCK, TERM_IN, TERM_OUT);
) BAF8 PS_AD_G_1 -8
) B2FE PS_VW_G_1 -2
) B2F6 PS_VW_G_1 -10
) B2F4 PS_VW_G_1 -12
) CB44004E0006E8FE INIT_PROCESS 68 78 6 -280
) END.
) D6 END_PROCESS
) CB2C000C0000B3FF INIT_PROCESS 44 12 0 -77
) FC HALT

```

```

E:
00 00 00 00 00 00 00 00 00 00 00 00 00 [ ..... ]
00 08 01 06 01 04 04 01 02 03 FF 00 00 49 [ .....I]
49 53 20 41 20 4E 45 57 20 44 41 59 00 [ IT IS A NEW DAY. ]

```

OK

	(01C1)	B0	FUNC_WORD				
	(01C2)	96FE01	PS_VW_L_1_FD	-2		1	
	(01C5)	91	PS_INDR				
	(01C6)	0A	CONSTANT	10			
	(01C7)	EA	DIV_WORD				
	(01C8)	CF8CFF	CALL_ROUTINE	-116			
	(01CB)	A3	COPY_BYTE				
	(01CC)	8AB6	PS_AD_G_1	-74			
	(01CE)	05	CONSTANT	5			
	(01CF)	A7	INDEX_11				
	(01D0)	B0	FUNC_WORD				
	(01D1)	96FE01	PS_VW_L_1_FD	-2		1	
	(01D4)	91	PS_INDR				
	(01D5)	0A	CONSTANT	10			
	(01D6)	EC	MOD				
	(01D7)	CF7DFF	CALL_ROUTINE	-131			
	(01DA)	A3	COPY_BYTE				
35:	(01DB)		BUFFER[6]:=',';				
	(01DB)	8AB6	PS_AD_G_1	-74			
	(01DD)	06	CONSTANT	6			
	(01DE)	A7	INDEX_11				
	(01DF)	9A3A	PS_CONST_1	58			
	(01E1)	A3	COPY_BYTE				
36:	(01E2)		BUFFER[7]:=ASCII(SEC DIV 10);			BUFFER[8]:=ASCII(SEC MOD 10	
	(01E2)	8AB6	PS_AD_G_1	-74			
	(01E4)	07	CONSTANT	7			
	(01E5)	A7	INDEX_11				
	(01E6)	B0	FUNC_WORD				
	(01E7)	96FE02	PS_VW_L_1_FD	-2		2	
	(01EA)	91	PS_INDR				
	(01EB)	0A	CONSTANT	10			
	(01EC)	EA	DIV_WORD				
	(01ED)	CF67FF	CALL_ROUTINE	-153			
	(01F0)	A3	COPY_BYTE				
	(01F1)	8AB6	PS_AD_G_1	-74			
	(01F3)	08	CONSTANT	8			
	(01F4)	A7	INDEX_11				
	(01F5)	B0	FUNC_WORD				
	(01F6)	96FE02	PS_VW_L_1_FD	-2		2	
	(01F9)	91	PS_INDR				
	(01FA)	0A	CONSTANT	10			
	(01FB)	EC	MOD				
	(01FC)	CF58FF	CALL_ROUTINE	-168			
	(01FF)	A3	COPY_BYTE				
37:	(0200)		BUFFER[9]:=NUL;				
	(0200)	8AB6	PS_AD_G_1	-74			
	(0202)	09	CONSTANT	9			
	(0203)	A7	INDEX_11				
	(0204)	A0	COPY_ZERO				
38:	(0205)		END;				
	(0205)	FE02	POP	2			
39:	(0207)		TERM_OUT.WRITE(BUFFER, NEWLINE);				
	(0207)	51	FPS_VW_G+06				
	(0208)	8AB6	PS_AD_G_1	-74			
	(020A)	01	CONSTANT	1			
	(020B)	CFF7FD	CALL_ROUTINE	-521			
40:	(020E)		END;				
	(020E)	442?	JUMP	34			
41:	(0210)	<>	BEGIN END;				
	(0210)	4420	JUMP	32			
42:	(0212)		END;				
	(0212)	CE470C	CASE_JUMP	71	12		
	(0215)	FAFF7AFFF6FFF4FF		-6	-134	-10	-12
	(021D)	F2FFF0FFEEFFECFF		-14	-16	-18	-20
	(0225)	EAF8FFE6FFE4FF		-22	-24	-26	-28

```

( 0175) 449C          JUMP          156
2: ( 0177) 'S' (SET TIME) ;
3: ( 0177)          BEGIN
7: ( 0177)          WITH T DO
( 0177) 8AB2          PS_AD_G_1      -78
0: ( 0179)          BEGIN
1: ( 179)          HOUR:=23;
( 0179) 86FE          PS_VW_L_1      -2
( 017B) 17          CONSTANT 23
( 017C) A3          COPY_BYTE
2: ( 017D)          MIN :=50;
( 017D) 96FE01       PS_VW_L_1_FD    -2      1
( 0180) 9A32        PS_CONST_1      50
( 0182) A3          COPY_BYTE
3: ( 0183)          SEC :=10;
( 0183) 96FE02       PS_VW_L_1_FD    -2      2
( 0186) 0A          CONSTANT 10
( 0187) A3          COPY_BYTE
4: ( 0188)          END;
( 0188) FE02        POP          2
5: ( 018A)          CLOCK.SETTIME(T);
( 018A) 53          FPS_VW_G+10
( 018B) 8AB2        PS_AD_G_1      -78
( 018D) CF7DFF      CALL_ROUTINE   -131
26: ( 0190)          END;
( 0190) 44A0        JUMP          160
27: ( 0192) 'G' (GET TIME) ;
28: ( 0192)          BEGIN
29: ( 0192)          CLOCK.GETTIME(T);
( 0192) 53          FPS_VW_G+10
( 0193) 8AB2        PS_AD_G_1      -78
( 0195) CF7FFF      CALL_ROUTINE   -129
30: ( 0198)          WITH T DO
( 0198) 8AB2        PS_AD_G_1      -78
31: ( 019A)          BEGIN
32: ( 019A)          BUFFER[1]:=ASCII(HOUR DIV 10); BUFFER[2]:=ASCII(HOUR MOD 1
( 019A) 8AB6        PS_AD_G_1      -74
( 019C) 01          CONSTANT 1
( 019D) A7          INDEX_11
( 019E) B0          FUNC_WORD
( 019F) 86FE        PS_VW_L_1      -2
( 01A1) 91          PS_INDB
( 01A2) 0A          CONSTANT 10
( 01A3) EA          DIV_WORD
( 01A4) CFB0FF      CALL_ROUTINE   -80
( 01A7) A3          COPY_BYTE
( 01A8) 8AB6        PS_AD_G_1      -74
( 01AA) 02          CONSTANT 2
( 01AB) A7          INDEX_11
( 01AC) B0          FUNC_WORD
( 01AD) 86FE        PS_VW_L_1      -2
( 01AF) 91          PS_INDB
( 01B0) 0A          CONSTANT 10
( 01B1) EC          MOD
( 01B2) CFA2FF      CALL_ROUTINE   -94
( 01B5) A3          COPY_BYTE
33: ( 01B6)          BUFFER[3]:=': ';
( 01B6) 8AB6        PS_AD_G_1      -74
( 01B8) 03          CONSTANT 3
( 01B9) A7          INDEX_11
( 01BA) 9A3A        PS_CONST_1      58
( 01BC) A3          COPY_BYTE
34: ( 01BD)          BUFFER[4]:=ASCII(MIN DIV 10); BUFFER[5]:=ASCII(MIN MOD 1
( 01BD) 8AB6        PS_AD_G_1      -74
( 01BE) 04          CONSTANT 4

```

```

03: ( 0141) (#####)
04: ( 0141) (#####)
05: ( 0141) # CLOCKPROCESS #
06: ( 0141) (#####)
07: ( 0141)
08: ( 0141) TYPE CLOCKPROCESS = PROCESS (PULSF: CLOCKPULSE;
09: ( 0141) CLOCK: CLOCK_MONITOR);
10: ( 0141)
11: ( 0141) BEGIN
12: ( 0141) CYCLE CLOCK.TICK; PULSE.TICK; END;
13: ( 0141) S1 FPS_VW_G+06
14: ( 0142) CF89FF CALL_ROUTINE -119
15: ( 0145) S2 FPS_VW_G+08
16: ( 0146) CFE8FF CALL_ROUTINE -21
17: ( 0149) 43F7 JUMP -9
18: ( 014B) END;
19: ( 014B) D6 END_PROCESS
20: ( 014C)
21: ( 014C) (#####)
22: ( 014C) # OPERATORPROCESS #
23: ( 014C) (#####)
24: ( 014C)
25: ( 014C) TYPE OPERATORPROCESS = PROCESS ( CLOCK: CLOCK_MONITOR;
26: ( 014C) TERM_IN: UART_READ;
27: ( 014C) TERM_OUT: UART_WRITE);
28: ( 014C)
29: ( 014C) VAR BUFFER: LINE;
30: ( 014C) LENGTH: INT;
31: ( 014C) T: TIME;
32: ( 014C)
33: ( 014C) FUNCTION NUMBER(CHARACTER: CHAR): INT;
34: ( 014C) BEGIN NUMBER:=ORD(CHARACTER) - ORD('0'); END;
35: ( 014C) C10A00 ENTER 10 0
36: ( 014F) 72 FPS_AD_L+12
37: ( 0150) 31 FPS_VW_L+10
38: ( 0151) 9A30 PS_CONST_1 48
39: ( 0153) E6 SUB_WORD
40: ( 0154) A2 COPY_WORD
41: ( 0155) D1 EXIT
42: ( 0156)
43: ( 0156) FUNCTION ASCII(NUMBER: INT): CHAR;
44: ( 0156) BEGIN ASCII:=CHR(NUMBER + ORD('0')); END;
45: ( 0156) C10A00 ENTER 10 0
46: ( 0159) 72 FPS_AD_L+12
47: ( 015A) 31 FPS_VW_L+10
48: ( 015B) 9A30 PS_CONST_1 48
49: ( 015D) E4 ADD_WORD
50: ( 015E) A2 COPY_WORD
51: ( 015F) D1 EXIT
52: ( 0160)
53: ( 0160) BEGIN
54: ( 0160) CYCLE
55: ( 0160) TERM_OUT.WRITE(NUL2, PROMPT);
56: ( 0160) S1 FPS_VW_G+06
57: ( 0161) 891A00 PS_AD_CN_2 26
58: ( 0164) 00 CONSTANT 0
59: ( 0165) CF9DFE CALL_ROUTINE -355
60: ( 0168) TERM_IN.READ(BUFFER, LENGTH);
61: ( 0168) S2 FPS_VW_G+08
62: ( 0169) 8AB6 PS_AD_G_1 -74
63: ( 016B) 8AB5 PS_AD_G_1 -75
64: ( 016D) CF08FF CALL_ROUTINE -248
65: ( 0170) CASE BUFFER[1] OF
66: ( 0170) 8AB6 PS_AD_G_1 -74
67: ( 0172) 01 CONSTANT 1
68: ( 0173) A7 INDEX 1.1

```

```

( 0106) CFFCFE          CALL_ROUTINE      -260
8: ( 0109)                END;
9: ( 0109)                END
0: ( 0109)                END
1: ( 0109)                END;
   ( 0109) FE02          POP                2
2: ( 010B)                END;
   ( 010B) D3           EXIT_MON
3: ( 010C)
4: ( 010C) PROCEDURE ENTRY SETTIME(T: TIME);
5: ( 010C) BEGIN
   ( 010C) C30C00        ENTER_MON         12      0
6: ( 010F)                CLOCKTIME := T;
   ( 010F) 8AFD         PS_AD_G_1         -3
   ( 0111) 31           FPS_VW_L+10
   ( 0112) A40300       COPY_STRUC         3
7: ( 0115)                END;
   ( 0115) D3           EXIT_MON
8: ( 0116)
9: ( 0116) PROCEDURE ENTRY GETTIME(VAR T: TIME);
0: ( 0116) BEGIN
   ( 0116) C30C00        ENTER_MON         12      0
1: ( 0119)                T := CLOCKTIME;
   ( 0119) 31           FPS_VW_L+10
   ( 011A) 8AFD         PS_AD_G_1         -3
   ( 011C) A40300       COPY_STRUC         3
2: ( 011F)                END;
   ( 011F) D3           EXIT_MON
3: ( 0120)
4: ( 0120) BEGIN
   ( 0120) C20A00        BEGIN_MON         10      0
5: ( 0123)                WITH CLOCKTIME DO BEGIN HOUR:=0; MIN:=0; SEC:=0; END;
   ( 0123) 8AFD         PS_AD_G_1         -3
   ( 0125) 86FE         PS_VW_L_1         -2
   ( 0127) A0           COPY_ZERO
   ( 0128) 96FE01       PS_VW_L_1_FD      -2      1
   ( 012B) A0           COPY_ZERO
   ( 012C) 96FE02       PS_VW_L_1_FD      -2      2
   ( 012F) A0           COPY_ZERO
   ( 0130) FE02        POP                2
6: ( 0132)                END;
   ( 0132) D2           END_MON
7: ( 0133)
8: ( 0133) (*****
9: ( 0133) * CLOCKPULSE *
0: ( 0133) *****
1: ( 0133)
2: ( 0133) TYPE CLOCKPULSE = DEVICE_MON (SELECTOR: INT);
3: ( 0133)
4: ( 0133) PROCEDURE ENTRY TICK;
5: ( 0133) BEGIN
   ( 0133) C30A00        ENTER_MON         10      0
6: ( 0136)                DOIO;
   ( 0136) DA           DOIO
7: ( 0137)                OUT(#00, CLOCK_WORD);
   ( 0137) 00          CONSTANT          0
   ( 0138) 94020B       PS_CONST_2        2050
   ( 013B) B7          OUT
8: ( 013C)                END;
   ( 013C) D3           EXIT_MON
9: ( 013D)
0: ( 013D) BEGIN
   ( 013D) C20A00        BEGIN_MON         10      0
1: ( 0140)                ( SET UP TO START THE CLOCK )
2: ( 0140)                END;

```

```

( 00BE) 8AFF PS_AD_G_1 -1
( 00C0) 940201 PS_CONST_2 258
( 00C3) B6 INN
( 00C4) A3 COPY_BYTE
3: ( 00C5) THROWAWAY:=INN(DATA_WORD);
( 00C5) 8AFF PS_AD_G_1 -1
( 00C7) 940201 PS_CONST_2 258
( 00CA) B6 INN
( 00CB) A3 COPY_BYTE
4: ( 00CC) END;
( 00CC) D2 END_MON
5: ( 00CD)
6: ( 00CD) (#####)
7: ( 00CD) # CLOCK_MONITOR #
8: ( 00CD) (#####)
9: ( 00CD)
0: ( 00CD) TYPE CLOCK_MONITOR = MONITOR (TERM_OUT: UART_WRITE);
1: ( 00CD) VAR CLOCKTIME: TIME;
2: ( 00CD)
3: ( 00CD) PROCEDURE ENTRY TICK;
4: ( 00CD) BEGIN
( 00CD) C30A00 ENTER_MON 10 0
5: ( 00D0) WITH CLOCKTIME DO
( 00D0) 8AFD PS_AD_G_1 -3
6: ( 00D2) BEGIN
7: ( 00D2) INC(SEC);
( 00D2) 96FE02 PS_VW_L_1_FD -2 2
( 00D5) 9F INC_BYTE
8: ( 00D6) IF SEC=60 THEN
( 00D6) 96FE02 PS_VW_L_1_FD -2 2
( 00D9) 91 PS_INDB
( 00DA) 9A3C PS_CONST_1 60
( 00DC) EE EQ_WORD
( 00DD) 4C2B FALSEJUMP 43
9: ( 00DF) BEGIN
10: ( 00DF) SEC:=0; INC(MIN);
( 00DF) 96FE02 PS_VW_L_1_FD -2 2
( 00E2) A0 COPY_ZERO
( 00E3) 96FE01 PS_VW_L_1_FD -2 1
( 00E6) 9F INC_BYTE
11: ( 00E7) IF MIN=60 THEN
( 00E7) 96FE01 PS_VW_L_1_FD -2 1
( 00EA) 91 PS_INDB
( 00EB) 9A3C PS_CONST_1 60
( 00ED) EE EQ_WORD
( 00EE) 4C1A FALSEJUMP 26
12: ( 00F0) BEGIN
13: ( 00F0) MIN:=0; INC(HOUR);
( 00F0) 96FE01 PS_VW_L_1_FD -2 1
( 00F3) A0 COPY_ZERO
( 00F4) 86FE PS_VW_L_1 -2
( 00F6) 9F INC_BYTE
14: ( 00F7) IF HOUR=24 THEN
( 00F7) 86FE PS_VW_L_1 -2
( 00F9) 91 PS_INDB
( 00FA) 18 CONSTANT 24
( 00FB) EE EQ_WORD
( 00FC) 4C0C FALSEJUMP 12
15: ( 00FE) BEGIN
16: ( 00FE) HOUR:=0;
( 00FE) 86FE PS_VW_L_1 -2
( 0100) A0 COPY_ZERO
17: ( 0101) TERM_OUT.WRITE('IT IS A NEW DAY(:0:)', NEWLINE);
( 0101) 51 FPS_VW_G+06
( 0102) B91C00 PS_AD_CN_2 28

```



```

( 007D) 91 PS_INDB
( 007E) 9A4A PS_CONST_1 74
( 0080) F0 LS_WORD
( 0081) 4C03 FALSEJUMP 3
( 0083) 31 FPS_VW_L+10
( 0084) 9F INC_BYTE
5: ( 0085) DOIO;
( 0085) DA DOIO
6: ( 0086) LASTCHAR := CHR(INN(DATA_WORD)); (GET A CHARACTER)
( 0086) 6F FPS_AD_L-01
( 0087) 940201 PS_CONST_2 258
( 008A) B6 INN
( 008B) A3 COPY_BYTE
7: ( 008C) IF LASTCHAR=BS THEN
( 008C) 2F FPS_VB_L-01
( 008D) 08 CONSTANT 8
( 008E) EE EQ_WORD
( 008F) 4C0D FALSEJUMP 13
8: ( 0091) BEGIN
9: ( 0091) IF LENGTH >= 2 THEN DEC(LENGTH);
( 0091) 31 FPS_VW_L+10
( 0092) 91 PS_INDB
( 0093) 02 CONSTANT 2
( 0094) F4 NL_WORD
( 0095) 4C03 FALSEJUMP 3
( 0097) 31 FPS_VW_L+10
( 0098) 9D DEC_BYTE
0: ( 0099) DEC(LENGTH);
( 0099) 31 FPS_VW_L+10
( 009A) 9D DEC_BYTE
1: ( 009B) END
2: ( 009B) ELSE MESSAGE[LENGTH] := LASTCHAR;
( 009B) 4407 JUMP 7
( 009D) 32 FPS_VW_L+12
( 009E) 31 FPS_VW_L+10
( 009F) 91 PS_INDB
( 00A0) A7 INDEX_11
( 00A1) 2F FPS_VB_L-01
( 00A2) A3 COPY_BYTE
3: ( 00A3) IF LASTCHAR <> CR THEN
( 00A3) 2F FPS_VB_L-01
( 00A4) 0D CONSTANT 13
( 00A5) F6 NE_WORD
( 00A6) 4C06 FALSEJUMP 6
4: ( 00A8) BEGIN
5: ( 00A8) TERM_OUT.ECHO(LASTCHAR);
( 00A8) 52 FPS_VW_G+08
( 00A9) 2F FPS_VB_L-01
( 00AA) CFACFF CALL_ROUTINE -84
6: ( 00AD) END;
7: ( 00AD) UNTIL (LASTCHAR = CR);
( 00AD) 2F FPS_VB_L-01
( 00AE) 0D CONSTANT 13
( 00AF) EE EQ_WORD
( 00B0) 4BCR FALSEJUMP -53
8: ( 00B2) TERM_OUT.WRITE(NUL2, NEWLINE);
( 00B2) 52 FPS_VW_G+08
( 00B3) 891A00 PS_AD_CN_2 26
( 00B6) 01 CONSTANT 1
( 00B7) CF4BFF CALL_ROUTINE -181
9: ( 00BA) END;
( 00BA) D3 EXIT_MON
0: ( 00BB)
1: ( 00BB) BEGIN
( 00BB) C20A00 BEGIN MON 10 0

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( 0041) EE EQ_WORD
( 0042) 4C08 FALSEJUMP 8
73: ( 0044) BEGIN OUT(ORD(')'), DATA_WORD); DOIO; END;
( 0044) 9A3E PS_CONST_1 62
( 0046) 940201 PS_CONST_2 258
( 0049) B7 OUT
( 004A) DA DOIO
74: ( 004B) OUT(#3D, CTRL_WORD); (TRANSMIT INHIBIT OTHERWISE SAME AS ABOVE)
( 004B) 9A3D PS_CONST_1 61
( 004D) 940301 PS_CONST_2 259
( 0050) B7 OUT
75: ( 0051) THROWAWAY := INN(CTRL_WORD);
( 0051) 6D FPS_AD_L-03
( 0052) 940301 PS_CONST_2 259
( 0055) B6 INN
( 0056) A2 COPY_WORD
76: ( 0057) END;
( 0057) D3 EXIT_MON
77: ( 0058)
78: ( 0058)
79: ( 0058) PROCEDURE ENTRY ECHO(CHRS:CHAR);
80: ( 0058) BEGIN
( 0058) C30C00 ENTER_MON 12 0
81: ( 005B) OUT(#BD, CTRL_WORD);
( 005B) 9ABD PS_CONST_1 189
( 005D) 940301 PS_CONST_2 259
( 0060) B7 OUT
82: ( 0061) OUT(ORD(CHRS), DATA_WORD);
( 0061) 31 FPS_VW_L+10
( 0062) 940201 PS_CONST_2 258
( 0065) B7 OUT
83: ( 0066) OUT(#3D, CTRL_WORD);
( 0066) 9A3D PS_CONST_1 61
( 0068) 940301 PS_CONST_2 259
( 006B) B7 OUT
84: ( 006C) END;
( 006C) D3 EXIT_MON
85: ( 006D)
86: ( 006D) BEGIN
( 006D) C20A00 BEGIN_MON 10 0
87: ( 0070) OUT(#3D, CTRL_WORD);
( 0070) 9A3D PS_CONST_1 61
( 0072) 940301 PS_CONST_2 259
( 0075) B7 OUT
88: ( 0076) END;
( 0076) D2 END_MON
89: ( 0077)
90: ( 0077) (*****
91: ( 0077) * UART_READ *
92: ( 0077) *****)
93: ( 0077)
94: ( 0077) TYPE UART_READ = DEVICE_MON (TERM_OUT; UART_WRITE;
95: ( 0077) SELECTOR; INT);
96: ( 0077)
97: ( 0077) VAR THROWAWAY: INT;
98: ( 0077)
99: ( 0077) PROCEDURE ENTRY READ( VAR MESSAGE: LINE; VAR LENGTH: INT);
100: ( 0077) VAR LASTCHAR: CHAR;
101: ( 0077) BEGIN
( 0077) C30E01 ENTER_MON 14 1
102: ( 007A) LENGTH:=0;
( 007A) 31 FPS_VW_L+10
( 007B) A0 COPY_ZERO
103: ( 007C) REPEAT
104: ( 007C) IF LENGTH<LINELENGTH THEN INC(LENGTH);

```

```

3: ( 0004) TYPE UART_WRITE = DEVICE MON (SELECTOR: INT);
4: ( 0004)
5: ( 0004) PROCEDURE ENTRY WRITE(MESSAGE: LINE; DISP: LINE_DISP);
6: ( 0004)   VAR I: INT;
7: ( 0004)     THROWAWAY: INTEGER;
8: ( 0004)   BEGIN
9: ( 0004) C30F03           ENTER_MON           14           3
10: ( J007)     I:=1;
11: ( 0007) 6F           FPS_AD_L-01
12: ( 0008) A1           COPY_ONE
13: ( 0009) OUT(#BD, CTRL_WORD) ( XMIT REQ., INT. EN., 8 DATA, 2 STOP, NO PA
14: ( 0009) 9ABD           PS_CONST_1           189
15: ( 000B) 940301       PS_CONST_2           259
16: ( 000E) B7           OUT
17: ( 000F) DOIO;
18: ( 000F) DA           DOIO
19: ( 0010) WHILE (MESSAGE[I] (>) NUL) AND (I<LINELENGTH) DO
20: ( 0010) 32           FPS_VW_L+12
21: ( 0011) 2F           FPS_VB_L-01
22: ( 0012) A7           INDEX_11
23: ( 0013) 91           PS_INDR
24: ( 0014) 00           CONSTANT 0
25: ( 0015) F6           NE_WORD
26: ( 0016) 2F           FPS_VB_L-01
27: ( 0017) 9A4A       PS_CONST_1           74
28: ( 0019) F0           LS_WORD
29: ( 001A) AD           AND
30: ( 001B) 4C0E       FALSEJUMP           14
31: ( 001D)   BEGIN
32: ( 001D)   OUT(ORD(MESSAGE[I]),DATA_WORD); (SEND A CHARACTER)
33: ( 001D) 32           FPS_VW_L+12
34: ( 001E) 2F           FPS_VB_L-01
35: ( 001F) A7           INDEX_11
36: ( 0020) 91           PS_INDR
37: ( 0021) 940201     PS_CONST_2           258
38: ( 0024) B7           OUT
39: ( 0025) DOIO;
40: ( 0025) DA           DOIO
41: ( 0026) INC(I);
42: ( 0026) 6F           FPS_AD_L-01
43: ( 0027) 9F           INC_BYTE
44: ( 0028) END;
45: ( 0028) 43F7       JUMP -25
46: ( 002A) IF (DISP=PROMPT) OR (DISP=NEWLINE) THEN
47: ( 002A) 31           FPS_VW_L+10
48: ( 002B) 00           CONSTANT 0
49: ( 002C) EE           EQ_WORD
50: ( 002D) 31           FPS_VW_L+10
51: ( 002E) 01           CONSTANT 1
52: ( 002F) EE           EQ_WORD
53: ( 0030) AE           OR
54: ( 0031) 4C0D       FALSEJUMP           13
55: ( 0033)   BEGIN OUT(ORD(CR), DATA_WORD); DOIO;
56: ( 0033) 0D           CONSTANT 13
57: ( 0034) 940201     PS_CONST_2           258
58: ( 0037) B7           OUT
59: ( 0038) DA           DOIO
60: ( 0039)   OUT(ORD(LF), DATA_WORD); DOIO;
61: ( 0039) 0A           CONSTANT 10
62: ( 003A) 940201     PS_CONST_2           258
63: ( 003D) B7           OUT
64: ( 003E) DA           DOIO
65: ( 003F)   END;
66: ( 003F) IF DISP=PROMPT THEN
67: ( 003F)   FPS_VW_L+10

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